

# The Art of Running Out of IPv6 Addresses

Benedikt Stockebrand  
Stepladder IT Training+Consulting GmbH

RIPE 77  
October 2018  
Amsterdam, The Netherlands

## NOTICE

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# About Me

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- Long term Unix, TCP/IP, Internet, architecture, operations, major cleanup guy
- Got serious with IPv6 in 2003
- Wrote a book about it in 2006
- Co-authored a study on secure IPv6 deployment for the German Federal Office for IT Security (BSI) in 2010
- Became one of your friendly IPv6 WG co-chairs in 2014
- Currently do a contract as “program manager” for IPv6 deployment at the datacenter operator for the state of Hessen (Germany)

# Part I

## Science

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**There is no shortage of addresses with IPv6!**



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- Think in bit utilization, not addresses.

## Part II

# Real World Strategies and Tactics



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- . . . but in doing so we achieve the inconceivable!

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- ... or 8 ~~bytes~~ octets



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- Once the effects of DS-Lite etc. wear off we can hope for reasonable predictions again.
- Rough estimate: In 50–100 years we will outgrow IPv6.



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- ... or even allocate at nibble boundaries.



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  - ...
- What about a Tech-C mail address?



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- Facts are for wimps.



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- ... and then give out /62s.
- Keep the rest “for future use”.

- It works with IPv4, why shouldn't it work with IPv6, too?







Picture: Wikipedia



Georg Wilhelm Friedrich Hegel

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Georg Wilhelm Friedrich Hegel  
German Philosopher

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Georg Wilhelm Friedrich Hegel  
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Picture: Wikipedia

Georg Wilhelm Friedrich Hegel  
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*The only thing we learn from history  
is that we learn nothing from history.*



Stepladder IT  
Training+Consulting GmbH  
Benedikt Stockebrand

Fichardstr. 38  
D-60322 Frankfurt/Main  
Germany

`contact@stepladder-it.com`

Web pages:

`http://www.stepladder-it.com/`

`http://www.benedikt-stockebrand.de/`

Video blog:

`http://www.stepladder-it.com/bivblog/`